

Outline

- 1 Formulating the Problem
- 2 Nearest Neighbors for Texts
- 3 Proving Hardness of Nearest Neighbors

Informal Problem Statement

To preprocess a database of n objects so that given a query object, one can effectively determine its nearest neighbors in database

Algorithms for Nearest Neighbors

Background and Two Challenges

Yury Lifshits

Steklov Institute of Mathematics at St.Petersburg

<http://logic.pdmi.ras.ru/~yura>

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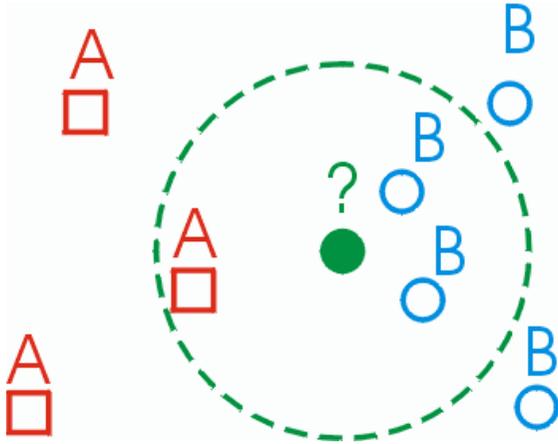
Part I Formulating the Problem

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First Application (1960s)

Nearest neighbors for classification:



Picture from <http://cgm.cs.mcgill.ca/~sooss/cs644/projects/perrier/Image25.gif>

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Applications

What applications of nearest neighbors do you know?

- Text classification
- Statistical data analysis, e.g. medicine diagnosis
- Pattern recognition: characters, faces
- Code plagiarism detection
- Coding theory
- Data compression
- **Web:** recommendation systems, on-line ads, personalized news aggregation, long queries in web search, near-duplicates detection

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Data Model in General

Formalization for nearest neighbors consists of:

- Representation format for objects
- Similarity function

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Basic Data Models (1/2)

- Vector Model
 - Similarity: l^2 , scalar product, cosine
- String Model
 - Similarity: Hamming distance, edit distance
- Black-box model
 - Similarity: given by oracle
 - The only knowledge is triangle inequality

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Basic Data Models (2/2)

- Set Model
 - Similarity: size of intersection
- Small graphs
 - Similarity: structure/labels matching

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Algorithmic Approaches to NN

- Divide and conquer
- Traversal techniques
- Look-up techniques
- Contractive and low-distortion embeddings
- Tournament algorithms

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Part II Nearest Neighbors for Texts

Sparse Vector Model

Database: points in R^d ,
every point has at most $k \ll d$ nonzero coordinates

Similarity: scalar product

Constraints:
 $poly(n + d)$ for preprocessing time,
 $poly(k) \cdot polylog(n + d)$ for query

Open Problem: solve NN for sparse vector model
within given constraints

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Inverted Index

Preprocessing:

For every term store a list of all documents in database with nonzero weight on it

Query processing:

Retrieve all points that have at least one common term with the query document;
Perform linear scan on them

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Rare-Point Method

Cheating: we will search only for neighbors that have at least one common rare feature with query object

Preprocessing:

For every rare feature store a list of all objects in database having it

Query processing:

Retrieve all points that have at least one common rare feature with the query object;
Perform linear scan on them

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Probabilistic Analysis in a Nutshell

- We define a probability distribution over databases
- We define probability distribution over query objects
- We construct a solution that is efficient/accurate with high probability over “random” input/query

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Zipf Model

- Terms t_1, \dots, t_m
- To generate a document we take every t_i with probability $\frac{1}{i}$
- Database is n independently chosen documents
- Query document has exactly one term in every interval $[e^i, e^{i+1}]$
- Similarity between documents is defined as the number of common terms

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Magic Level Theorem

Magic Level $q = \sqrt{2 \log_e n}$

Theorem (Hoffmann, Lifshits and Nowotka, CSR'07)

- 1 With very high probability there exists a document in database having $q - \varepsilon$ **top** terms of query document
- 2 With very small probability there exists a document in database having **any** $q + \varepsilon$ overlap with query document

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Inclusions with Preprocessing (1/2)

Input

Family \mathcal{F} of subsets of U

Query task

Given a set $f_{new} \subseteq U$ to decide whether $\exists f \in \mathcal{F} : f_{new} \subseteq f$

Constraints

Data storage after preprocessing $poly(|\mathcal{F}| + |U|)$
Time for query processing $poly(|U|)$

Open problem: is there an algorithm satisfying given constraints?

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Part III Proving Hardness of Nearest Neighbors

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Inclusions with Preprocessing (2/2)

Reformulation in SAT style:

Input

DNF formula \mathcal{F} on n variables, without negations

Query task

Given an assignment x to evaluate $\mathcal{F}(x)$

Constraints

Data storage after preprocessing $poly(|\mathcal{F}|)$
Time for query processing $poly(n)$

Open problem: is there an algorithm satisfying given constraints?

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“NP Analogue” for Search Problems

Every problem in **SEARCH class** is characterized by poly-time computable Turing Machine M :

Input

Strings x_1, \dots, x_n , $|x_i| = m$

Query task

Given string y of length m to answer whether $\exists i : M(x_i, y) = \text{yes}$

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Tractable problems in SEARCH

Input

Strings x_1, \dots, x_n , $|x_i| = m$

Query task

Given string y of length m to answer whether $\exists i : M(x_i, y) = \text{yes}$

Tractable solution

Preprocessing in $\text{poly}(m, n)$ space

Query processing in $\text{poly}(m, \log n)$ time with RAM access to preprocessed database

Inclusions is in SEARCH. Is it tractable?

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Complete problems in SEARCH (1/2)

Program Search problem:

Input

Turing machines P_1, \dots, P_n

Query task

Given string y of length m to answer whether $\exists i : P_i(y) = \text{yes}$ after at most m steps

Open problem: is Program Search tractable?

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Complete problems in SEARCH (2/2)

Parallel Run problem:

Input

x_1, \dots, x_n

Query task

Given poly-time computable P to answer whether $\exists i : P(x_i) = \text{yes}$

Open problem: is Parallel Run tractable?

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Conclusions

- Any relevant work?
- How to improve this talk for the next time?
- **Give my open problems to your friends!**

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Summary

- Nearest neighbors for texts can be modelled by Euclidean space with sparse vectors
- No exact algorithms for NN in texts are known so far
- New approach to lower bounds for NN: SEARCH class and its complete problems

Thanks for your attention! Questions?

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References (1/2)

Search “**Lifshits**” or visit <http://logic.pdmi.ras.ru/~yura/>



B. Hoffmann, Y. Lifshits and D. Nowotka

Maximal Intersection Queries in Randomized Graph Models

<http://logic.pdmi.ras.ru/~yura/en/maxint-draft.pdf>



P.N. Yianilos

Data structures and algorithms for nearest neighbor search in general metric spaces

<http://www.pnylab.com/pny/papers/vptree/vptree.ps>



J. Zobel and A. Moffat

Inverted files for text search engines

<http://www.cs.mu.oz.au/~alastair/abstracts/zm06compsurv.html>



K. Teknomo

Links to nearest neighbors implementations

<http://people.revoledu.com/kardi/tutorial/KNN/resources.html>

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References (2/2)



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Two Algorithms for Nearest-Neighbor Search in High Dimensions

<http://www.ece.tuc.gr/~vsam/csalgo/kleinberg-stoc97-nn.ps>



P. Indyk and R. Motwani

Approximate nearest neighbors: towards removing the curse of dimensionality

<http://theory.csail.mit.edu/~indyk/nndraft.ps>



A. Andoni and P. Indyk

Near-Optimal Hashing Algorithms for Approximate Nearest Neighbor in High Dimensions

<http://theory.lcs.mit.edu/~indyk/FOCS06final.ps>



P. Indyk

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