

Four Results of Jon Kleinberg

A Talk for St.Petersburg Mathematical Society

Yury Lifshits

Steklov Institute of Mathematics at St.Petersburg

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Part I

History of Nevanlinna Prize

Career of Jon Kleinberg

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Outline

- 1 Nevanlinna Prize for Jon Kleinberg
 - History of Nevanlinna Prize
 - Who is Jon Kleinberg
- 2 Hubs and Authorities
- 3 Nearest Neighbors: Faster Than Brute Force
- 4 Navigation in a Small World
- 5 Bursty Structure in Streams

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Nevanlinna Prize

The Rolf Nevanlinna Prize is awarded **once every 4 years** at the International Congress of Mathematicians, for outstanding contributions in **Mathematical Aspects of Information Sciences** including:

- 1 All mathematical aspects of computer science, including complexity theory, logic of programming languages, analysis of algorithms, cryptography, computer vision, pattern recognition, information processing and modelling of intelligence.
- 2 Scientific computing and numerical analysis. Computational aspects of optimization and control theory. Computer algebra.

Only scientists **under 40** are eligible

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Previous Winners

- 1982 Robert Tarjan: data structures, graph algorithms
- 1986 Leslie Valiant: learning theory, complexity, parallel computing
- 1990 Alexander Razborov: work around P vs. NP
- 1994 Avi Wigderson: complexity and cryptography
- 1998 Peter Shor: quantum algorithm for factoring problem
- 2002 Madhu Sudan: coding theory, probabilistically checkable proofs and inapproximability

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Short Bio of Jon Kleinberg



- 1971 Jon Kleinberg was born in Boston
- 1993 Bachelor degree from Cornell
- 1996 Ph.D. from MIT (advisor Michel X. Goemans)
- Since 1996 Cornell faculty
- 2006 Nevanlinna Prize

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More about Jon Kleinberg

- According to DBLP: 108 papers and 85 coauthors for 1992-2006
- H-Index = 36 (according to scholar.google.com)
- Book “Algorithm Design” (2005, with Éva Tardos)
- NSF Career Award, ONR Young Investigator Award, MacArthur Foundation Fellowship, Packard Foundation Fellowship, Sloan Foundation Fellowship, “Faculty of the Year” Cornell’2002
- Strong connections to IBM Almaden Research Center
- Courses “The Structure of Information Networks” and “Randomized and High-Dimensional Algorithms”
- Chair of STOC’06

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Research Style of Jon Kleinberg

- **Direction:** from practical problems to mathematical ideas
- **Motivation:** make life better
- **Validation:** mathematical proofs **and** experiments
- **Connections with:** sociology
- **Key component:** new models/formalizations, not proofs

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Part II

Authoritative sources in a hyperlinked environment
Jon Kleinberg — SODA'98

2580 citations
according to scholar.google.com, May 2007

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Web Search: Formal Settings

- Every webpage is represented as a weighted set of keywords
- There are hyperlinks (directed edges) between webpages

Conceptual problem: define a relevance rank based on keyword weights and link structure of the web

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Challenge

How to define the most relevant webpage to “Bill Gates”?

Naive ideas

- By frequency of query words in a webpage
- By number of links from other **relevant** pages

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HITS Algorithm

- 1 Given a query construct a **focused subgraph** $F(query)$ of the web
- 2 Compute **hubs and authorities** ranks for all vertices in $F(query)$

Focused subgraph: pages with highest weights of query words **and** pages hyperlinked with them

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Hubs and Authorities

Mutual reinforcing relationship:

- A good **hub** is a webpage with many links **to** query-authoritative pages
- A good **authority** is a webpage with many links **from** query-related hubs

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Hubs and Authorities: Equations

$$a(p) \sim \sum_{q:(q,p) \in E} h(q)$$

$$h(p) \sim \sum_{q:(p,q) \in E} a(q)$$

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Hubs and Authorities: Solution

Initial estimate:

$$\forall p : a_0(p) = 1, h_0(p) = 1$$

Iteration:

$$a_{k+1}(p) = \sum_{q:(q,p) \in E} h_k(q)$$

$$h_{k+1}(p) = \sum_{q:(p,q) \in E} a_k(q)$$

We normalize \bar{a}_k, \bar{h}_k after every step

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Convergence Theorem

Theorem

Let M be the adjacency matrix of focused subgraph $F(\text{query})$. Then \bar{a}_k converges to principal eigenvector of $M^T M$ and \bar{h}_k converges to principal eigenvector of MM^T

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Lessons from Hubs and Authorities

- Link structure is useful for relevance sorting
- Link popularity is defined by linear equations
- Solution can be computed by iterative algorithm

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Part III

Two algorithms for nearest-neighbor search
in high dimensions
Jon Kleinberg — STOC'97

173 citations
according to scholar.google.com, May 2007

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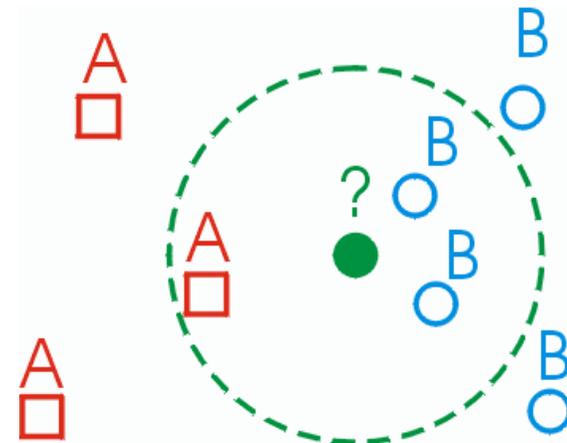
Informal Problem Statement

To preprocess a database of n objects
so that given a query object,
one can effectively determine
its nearest neighbors in database

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First Application (1960s)

Nearest neighbors for classification:



Picture from <http://cgm.cs.mcgill.ca/~sooss/cs644/projects/perrier/Image25.gif>

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Applications

What applications of nearest neighbors do you know?

- Statistical data analysis, e.g. medicine diagnosis
- Pattern recognition, e.g. for handwriting
- Code plagiarism detection
- Coding theory
- Future applications: recommendation systems, ads distribution, personalized news aggregation

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Approximate Nearest Neighbors

Definition

p is ε -approximate nearest neighbor for q
iff $\forall p' \in DB : d(p, q) \leq (1 + \varepsilon)d(p', q)$

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Challenge

Brute force algorithm

No preprocessing

$\mathcal{O}(nd)$ query time

for n points in d -dimensional space

Open Problem: Is there any preprocessing method with data structure of $\text{poly}(n + d)$ size and $o(nd)$ query time?

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Kleinberg Algorithm

Theorem

For every ε, δ there exists a data structure with $\mathcal{O}^*(d^2n)$ construction time and $\mathcal{O}(n + d \log^3 n)$ query processing time. It correctly answer ε -nearest neighbor queries with probability $1 - \delta$.

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Data Structure Construction

- 1 Choose $l = d \log^2 n \log^2 d$ random vectors $V = \{v_1, \dots, v_l\}$ with unit norm
- 2 Precompute all scalar products between database points and vectors from V

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Random Projection Test

Input: points x, y and q , vectors u_1, \dots, u_k

Question: what is smaller $|x - q|$ or $|y - q|$?

Test:

For all i compare $(x \cdot v_i - q \cdot v_i)$ with $(y \cdot v_i - q \cdot v_i)$
Return the point which has “smaller”
on majority of vectors

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Query Processing

- 1 Choose a random subset Γ of V , $|\Gamma| = \log^3 n$
- 2 Compute scalar products between query point q and vectors from Γ
- 3 Make a tournament for choosing a nearest neighbor:
 - 1 Draw a binary tree of height $\log n$
 - 2 Assign all database points to leaves
 - 3 For every internal point (say, x vs. y) make a random projection test using some vectors from Γ

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Part IV

The small-world phenomenon:
An algorithmic perspective
Jon Kleinberg — STOC'00

433 citations
according to scholar.google.com, May 2007

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Milgram's Small World Experiment

- 1 Starting point: Wichita/Omaha, endpoint: Boston
- 2 **Basic information** about a target contact person in Boston was initially sent to randomly selected individuals.
- 3 If recipient knew the contact person, he/she should **forward** the letter **directly** to that person
- 4 If recipient did not personally know the target then he/she should **forward** the package **to a friend** or relative they know personally that is more likely to know the target
- 5 When and if the package eventually reached the contact person in Boston, the researchers count the number of times it had been forwarded from person to person.

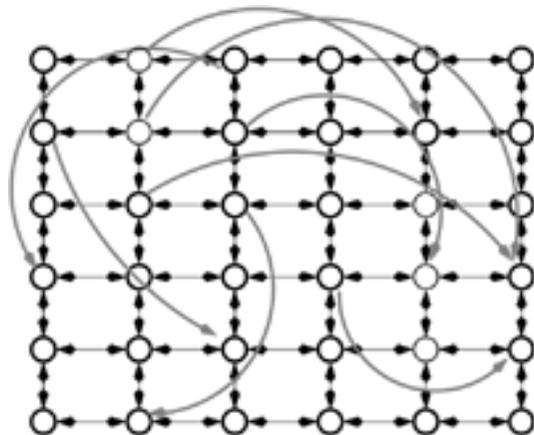
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Small World Model

- $n \times n$ grid of n^2 nodes
- Every node p is connected to its direct neighbors: right, left, up and down
- Additionally, every node p has an arc to a “random” node q , where probability for q to be chosen is proportional to $|p - q|^{-\alpha}$, $\alpha \geq 0$

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Small World Model



Picture from www.math.cornell.edu/~durrett/smw/kleinberg2.gif

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Navigability

A graph is **navigable**, if there exists **decentralized** algorithm finding connecting paths in $\text{polylog}(n)$ time

Whether small world is navigable?

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Kleinberg's Results

Theorem

For $\alpha = 2$ small world is navigable, for all other nonnegative values of α it is not.

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Streams and Bursts

- A stream of events
- Every event = set of keywords + time stamp

How should we identify time intervals with unusually high frequency of a specific keyword?

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Part V

Bursty and Hierarchical Structure in Streams
Jon Kleinberg — KDD'02

150 citations
according to scholar.google.com, May 2007

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Conceptual Solution

Hidden Markov Model methodology:

- There is a “creature” who generates our stream
- This creature can be described as a finite automaton of known structure but with unknown state sequence
- We will find “the most fitting” sequence of states for our data
- Based on this sequence we can identify all bursts

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Very Simple Example (1/2)

Keyword: “grant”

Events: every day either there is an email with this keyword or there is not

Example Data: we have email archive for two weeks

01110100001000

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Very Simple Example (1/2)

01110100001000

Automaton: two states “grant deadline” and “vacations”

Fitting function: 1 point penalty for mismatches “grant deadline — no grant emails” and “vacations — email with grants”, 1 point penalty for switching state of automaton

Optimal sequence: VDDDDDV VVVVVVV

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Algorithm for Detecting Bursts

How to compute the optimal state sequence?

Dynamic programming:

- For every day d and every state s we will compute the optimal state sequence for period $[1..d]$ ending with state s
- When a data for new day comes we try all values for yesterday and choose the best one
- For optimal sequence for the whole interval $[1..D]$ we just take the maximum over all states

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Home problem

Find an anagram for “KLEINBERG”

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Highlights

- Hubs and Authorities is an iterative algorithm for computing relevance rank
- Small world always can have small diameter but no decentralized method for finding short paths
- Bursts can be identified as states of imaginary automaton that generates event stream
- Nearest neighbors can be found by looking at projections to random vectors

Thank you for your attention!
Questions?

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Relevant Links

-  Official site of Nevanlinna Prize
<http://www.mathunion.org/Prizes/Nevanlinna/index.html>
-  Homepage of Jon Kleinberg
<http://www.cs.cornell.edu/home/kleinber/>
-  Jon Kleinberg at DBLP
<http://www.informatik.uni-trier.de/~ley/db/indices/a-tree/k/Kleinberg:Jon.M-.html>
-  IMU news release
<http://www.mathunion.org/medals/2006/kleinbergENG.pdf>
-  Interview of Jon Kleinberg to "Technology Research News"
http://www.trnmag.com/Stories/2005/120505/View_Jon_Kleinberg_120505.html
-  A talk by Jon Kleinberg on Yahoo! Video
<http://video.yahoo.com/video/play?vid=62055>

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References

All materials of this talk will be published at **my homepage**:
<http://logic.pdmi.ras.ru/~yura>



Jon Kleinberg

Authoritative sources in a hyperlinked environment — SODA'98

<http://www.cs.cornell.edu/home/kleinber/auth.pdf>



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The small-world phenomenon: An algorithmic perspective — STOC'00

<http://www.cs.cornell.edu/home/kleinber/swn.ps>



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Bursty and Hierarchical Structure in Streams — KDD'02

<http://www.cs.cornell.edu/home/kleinber/bhs.ps>



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Two algorithms for nearest-neighbor search in high dimensions — STOC'97

<http://www.cs.cornell.edu/home/kleinber/stoc97-nn.pdf>

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